

Maximizing a Tree Series in the Representation Space

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Overview

- 1 Starting Point: Metropolis Procedural Modeling
- 2 Problem Formulation
- 3 Working in the Representation Space
- 4 Experiments
- 5 Conclusion

Overview

1 Starting Point: Metropolis Procedural Modeling

2 Problem Formulation

- Preliminaries: Tree Series and the Representation Space
- Motivations and Problematic

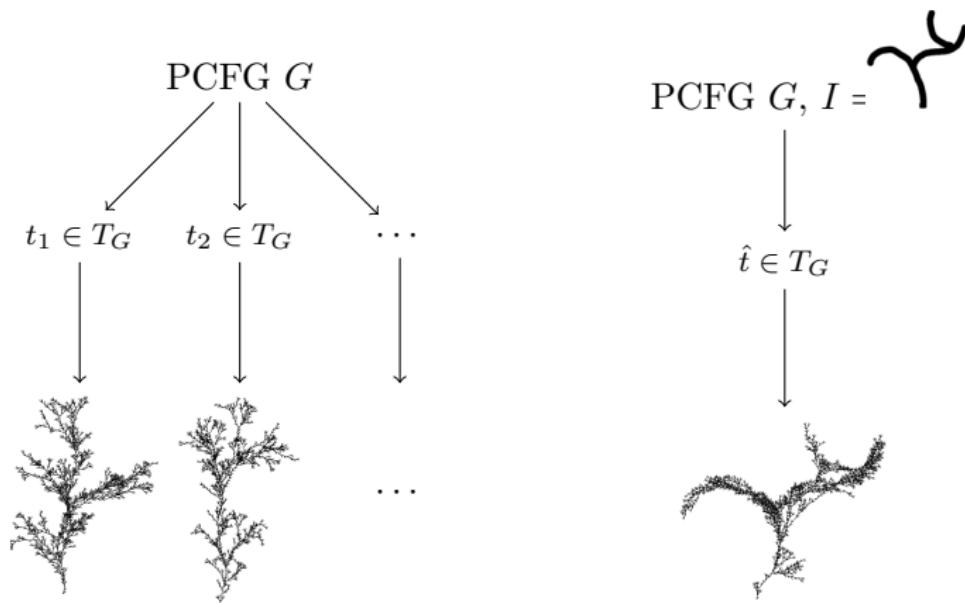
3 Working in the Representation Space

- Complexity Study
- Metropolis-Hastings in the Representation Space

4 Experiments

5 Conclusion

Metropolis Procedural Modeling [Talton et al., 2011]



2 steps:

- Define a *posterior* distribution $p(t|I) \propto \pi(t)L(I|t)$ on T_G
- Find $\hat{t} \in T_G$ maximizing $p(\cdot|I) \Rightarrow$ Metropolis-Hastings

Metropolis-Hastings Algorithm

- $p : \mathcal{X} \rightarrow \mathbb{R}_+$ such that $Z = \int_{\mathcal{X}} p(x) dx < \infty$
- $\Rightarrow \hat{p} : x \mapsto p(x) / Z$ is a probability distribution on \mathcal{X} .

Metropolis-Hastings Algorithm

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⇒ $\hat{p} : x \mapsto p(x)/Z$ is a probability distribution on \mathcal{X} .

To sample from \hat{p} :

- (i) Choose a *jump distribution* $q_x(\cdot)$ (distribution on \mathcal{X} for each $x \in \mathcal{X}$).
- (ii) Build a Markov chain in \mathcal{X} :

Input : $x_n \in \mathcal{X}$

Returns : $x_{n+1} \in \mathcal{X}$

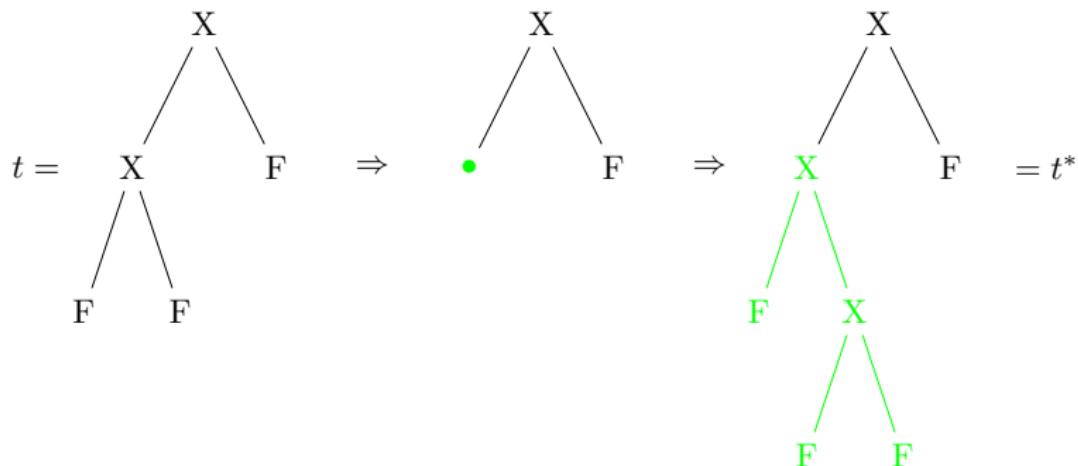
1: Draw a candidate $x^* \sim q_{x_n}(\cdot)$

2: Accept x^* (i.e. $x_{n+1} \leftarrow x^*$, otherwise $x_{n+1} \leftarrow x_n$) with probability

$$\alpha(x_n, x^*) = \min \left\{ 1, \frac{p(x^*) q_{x^*}(x_n)}{p(x_n) q_{x_n}(x^*)} \right\}$$

Maximizing $p(\cdot|I)$ with Metropolis-Hastings

- Just need to choose a jump distribution $q_t(\cdot)$:



- And build the Markov chain in T_G :

$$t_0 \xrightarrow{MH} t_1 \xrightarrow{MH} t_2 \xrightarrow{MH} t_3 \xrightarrow{MH} \dots$$

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Preliminaries: Tree Series

- Set of *tree series* : $\mathbb{R}\langle\langle \mathcal{F} \rangle\rangle = \{\phi : T_{\mathcal{F}} \rightarrow \mathbb{R}\}$
- $r \in \mathbb{R}\langle\langle \mathcal{F} \rangle\rangle$ is *rational* $\Leftrightarrow r$ is recognizable by a WTA

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Theorem

A tree series $r \in \mathbb{R}\langle\langle \mathcal{F} \rangle\rangle$ is rational if and only if it has a (finite dimensional) linear representation, i.e. there exist (V, μ, λ) s.t. :

- V is a finite dimensional vector space
- $\mu : T_{\mathcal{F}} \rightarrow V$ is a linear mapping
- $\lambda : V \rightarrow \mathbb{R}$ is a linear form s.t. $\forall t \in T_{\mathcal{F}} : r(t) = \lambda(\mu(t))$

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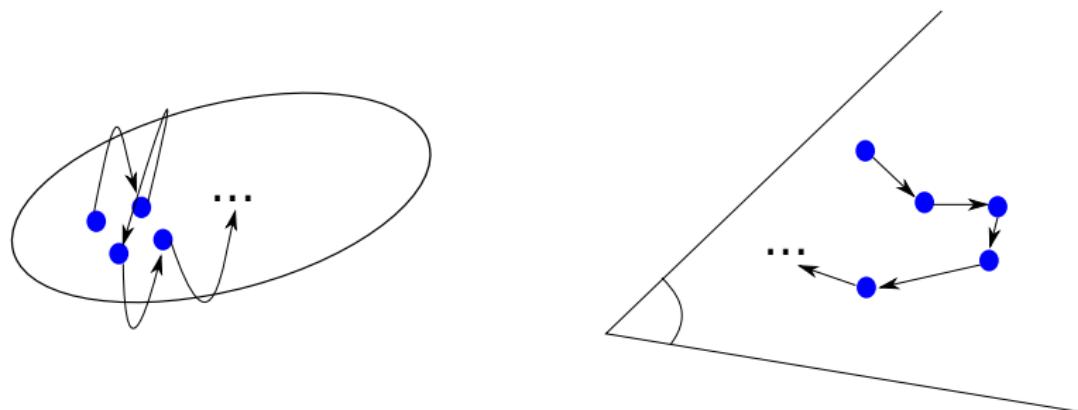
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- V is a finite dimensional vector space
- $\mu : T_{\mathcal{F}} \rightarrow V$ is a linear mapping
- $\lambda : V \rightarrow \mathbb{R}$ is a linear form s.t. $\forall t \in T_{\mathcal{F}} : r(t) = \lambda(\mu(t))$
- (V, μ) is a linear representation of $T_{\mathcal{F}}$.
- V is a representation space of $T_{\mathcal{F}}$.

Problematic

- (Metropolis Procedural Modeling [Talton et al., 2011])
 - ▶ Goal: find the tree maximizing $p(\cdot|I) \propto \pi(\cdot)L(I|\cdot)$.
 - ▶ Can we use the representation space induced by π ?



- Given a positive tree series $\phi : T_{\mathcal{F}} \rightarrow \mathbb{R}$, can we use a representation space of $T_{\mathcal{F}}$ to find

$$\hat{t} = \arg \max_{t \in T_{\mathcal{F}}} \phi(t)$$

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Complexity Study (1)

Definition (Max-RTS)

Input: A positive rational series $\phi \in \mathbb{R}\langle\langle \mathcal{F} \rangle\rangle$ and $\gamma \in \mathbb{Q}$.

Question: Is there a tree t such that $\phi(t) \geq \gamma$?

Theorem

- *Max-RTS is undecidable.*
- *With the added constraint that the support of ϕ is finite, Max-RTS is NP-hard.*

Complexity Study (2)

Definition (Ball-RTS)

Input : A linear representation $(\mathbb{R}^n, \mu, \lambda)$ of a positive rational series ϕ , a point $\mathbf{x} \in \mathbb{R}^n$ and $\gamma \in \mathbb{Q}$.

Question : Is there a $t \in \text{supp}(\phi)$ s.t. $\|\mu(t) - \mathbf{x}\| \leq \gamma$ (resp. $<$, $>$, \geq) ?

Theorem

- Ball-RTS is undecidable.
- With the added constraint that the support of ϕ is finite, Ball-RTS is NP-hard.

Complexity Study (2)

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Theorem

- Ball-RTS is undecidable.
- With the added constraint that the support of ϕ is finite, Ball-RTS is NP-hard.

Theorem

Given a linear representation (V, μ) of $T_{\mathcal{F}}$, the injectivity of μ is undecidable.

Maximization with Metropolis-Hastings in $T_{\mathcal{F}}$: MHTF

- Let $\phi : T_{\mathcal{F}} \rightarrow \mathbb{R}$ be the positive series to maximize and
 $\hat{\phi} : t \mapsto \phi(t)/Z$

To sample from $\hat{\phi}(\cdot)$:

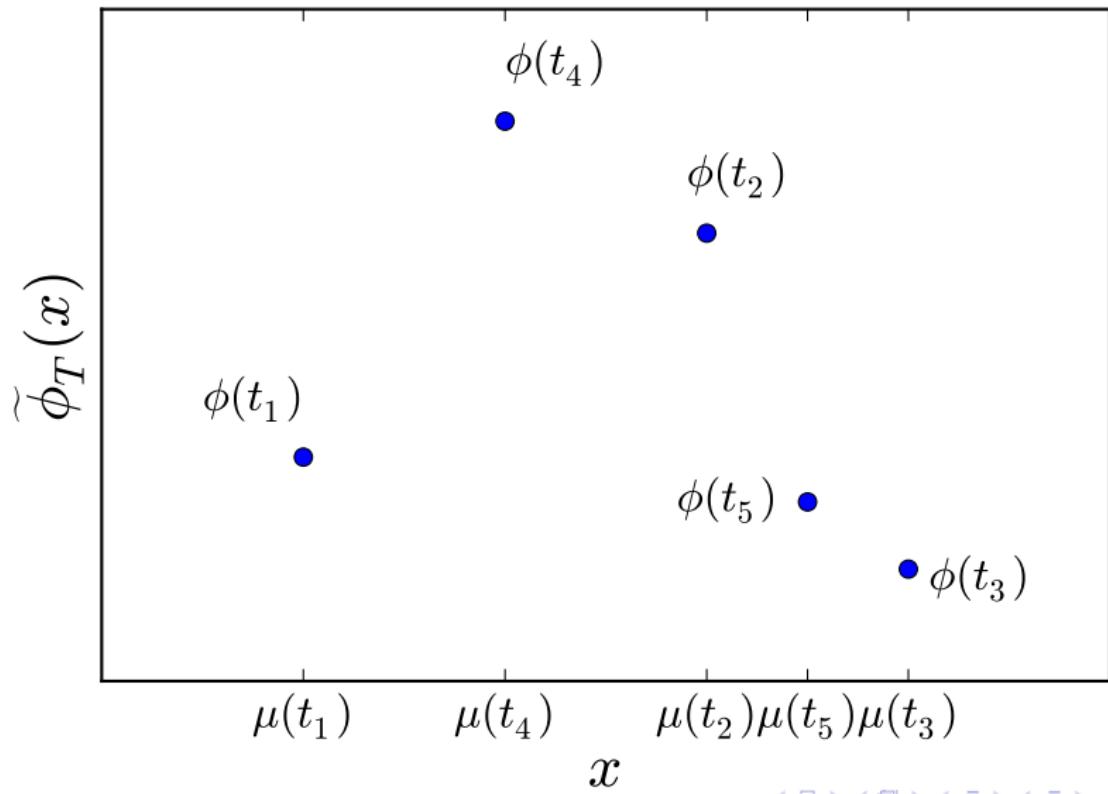
- Choose a distribution π on $T_{\mathcal{F}}$ and a distribution s_t on $C_{\mathcal{F}}(t)$.
- Build a Markov chain in $T_{\mathcal{F}}$:

Input : $t_n \in T_{\mathcal{F}}$

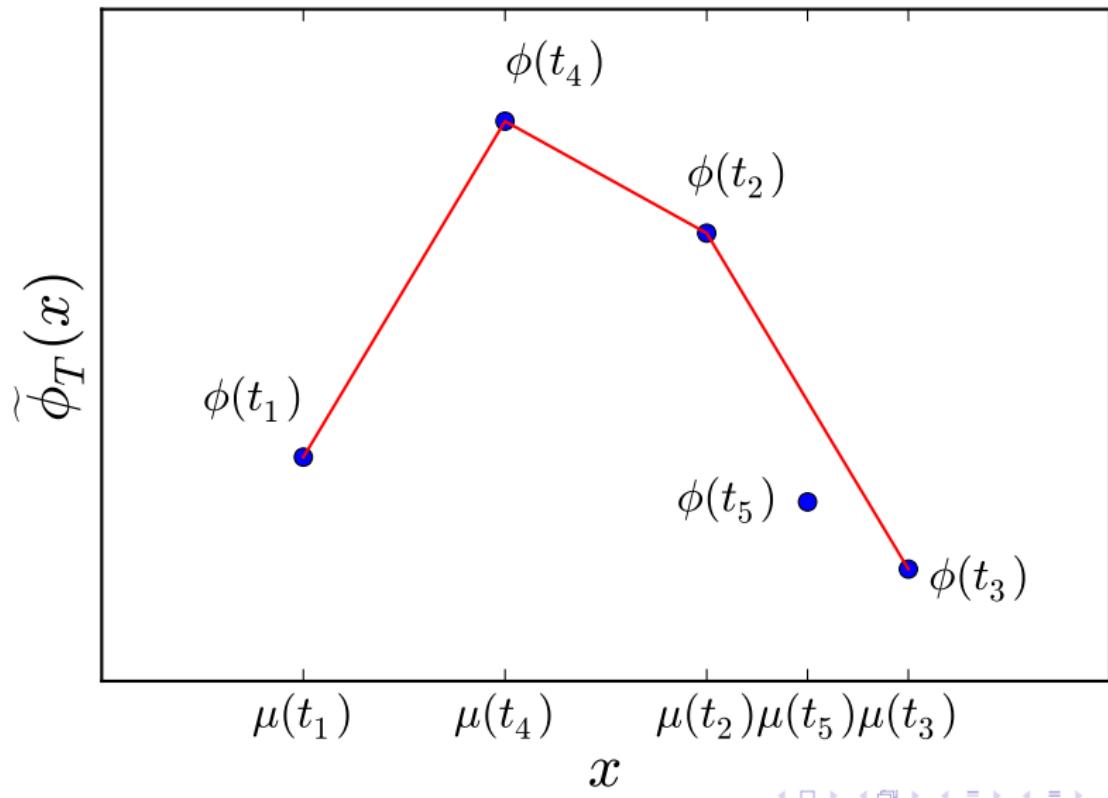
Returns : $t_{n+1} \in T_{\mathcal{F}}$

- Draw a context $c \in C_{\mathcal{F}}(t_n) \sim s_{t_n}(\cdot)$, and a subtree $\tau \in T_{\mathcal{F}} \sim c^{-1}\pi(\cdot)$
 - $t^* \leftarrow c[\tau]$
 - Accept t^* with probability $\alpha(t_n, t^*) = \min \left\{ 1, \frac{\phi(t^*) s_{t^*}(c) \pi(t_n)}{\phi(t_n) s_{t_n}(c) \pi(t^*)} \right\}$
-

Extending ϕ to the representation space



Extending ϕ to the representation space



Extending ϕ to the representation space

Let $\phi : T_{\mathcal{F}} \rightarrow \mathbb{R}$ be the series to maximize and (V, μ) a linear representation of $T_{\mathcal{F}}$.

Definition

Let $T \subseteq T_{\mathcal{F}}$ be non-empty. Define $\tilde{\phi}_T$ on $\mathcal{X} = \text{conv}(\mu(T))$ by

$$\tilde{\phi}_T(\mathbf{x}) = \sup_{\substack{n > 0, \alpha \in [0,1]^n \\ \sum \alpha_i = 1, t_1 \dots t_n \in T}} \left\{ \sum_{i=1}^n \alpha_i \phi(t_i) : \mathbf{x} = \sum_{i=1}^n \alpha_i \mu(t_i) \right\}$$

for all $\mathbf{x} \in \mathcal{X}$.

Theorem

$\tilde{\phi}_T$ is a well defined function, continuous on \mathcal{X} , and takes its maximum on $\mu(t_{\max})$.

Metropolis-Hastings for the distribution $\hat{\phi}_T$ in V

Input : $x_n \in \mathcal{X}$

Returns : $x_{n+1} \in \mathcal{X}$

- 1: Draw a candidate $x^* \sim q_{x_n}(\cdot)$
- 2: Accept x^* with probability

$$\alpha(x_n, x^*) = \min \left\{ 1, \frac{\hat{\phi}_T(x^*) q_{x^*}(x_n)}{\hat{\phi}_T(x_n) q_{x_n}(x^*)} \right\} = \min \left\{ 1, \frac{\tilde{\phi}_T(x^*)}{\tilde{\phi}_T(x_n)} \right\}$$

where $q_x(\cdot)$ is a symmetric and positive jump distribution.

Metropolis-Hastings for the distribution $\hat{\phi}_T$ in V

Input : $x_n \in \mathcal{X}$

Returns : $x_{n+1} \in \mathcal{X}$

- 1: Draw a candidate $x^* \sim q_{x_n}(\cdot)$
- 2: Accept x^* with probability

$$\alpha(x_n, x^*) = \min \left\{ 1, \frac{\hat{\phi}_T(x^*) q_{x^*}(x_n)}{\hat{\phi}_T(x_n) q_{x_n}(x^*)} \right\} = \min \left\{ 1, \frac{\tilde{\phi}_T(x^*)}{\tilde{\phi}_T(x_n)} \right\}$$

where $q_x(\cdot)$ is a symmetric and positive jump distribution.

Goal : Sample from $\hat{\phi}_{T_F}$

Idea : Build successive sets of trees $T_1 \subseteq T_2 \subseteq \dots$ while exploring \mathcal{X} .

Adaptive Metropolis-Hastings in V : MHV

Input : $\mathbf{x}_n \in \mathcal{X}$, $T_n = \{t_1, \dots, t_n\} \subseteq T_{\mathcal{F}}$

Returns : $\mathbf{x}_{n+1} \in \mathcal{X}$, $T_{n+1} \subseteq T_{\mathcal{F}}$

- 1: Draw $\mathbf{x}^* \sim q_{\mathbf{x}_n}(\cdot)$
- 2: Solve $\tilde{\phi}_{T_n}(\mathbf{x}^*) : \mathbf{x}^* = \sum_{i=1}^n \alpha_i \mu(t_i)$
- 3: Draw a tree $t \in T_n$ ($\sim p_{\alpha}$), a context $c \in C_{\mathcal{F}}(t)$ ($\sim q_t$), and a subtree $\tau \in T_{\mathcal{F}} \sim c^{-1}\pi(\cdot)$
- 4: $t_{n+1} \leftarrow c[\tau]$, $T_{n+1} \leftarrow T_n \cup \{t_{n+1}\}$
- 5: Accept \mathbf{x}^* with probability

$$\alpha(\mathbf{x}_n, \mathbf{x}^*) = \min \left\{ 1, \frac{\hat{\phi}_{T_n}(\mathbf{x}^*) q_{\mathbf{x}^*}(\mathbf{x}_n)}{\hat{\phi}_{T_n}(\mathbf{x}_n) q_{\mathbf{x}_n}(\mathbf{x}^*)} \right\} = \min \left\{ 1, \frac{\tilde{\phi}_{T_n}(\mathbf{x}^*)}{\tilde{\phi}_{T_n}(\mathbf{x}_n)} \right\}$$

where $q(\cdot, \cdot)$ is a symmetric and positive jump distribution.

Adaptive Metropolis-Hastings in V : MHV

Input : $\mathbf{x}_n \in \mathcal{X}$, $T_n = \{t_1, \dots, t_n\} \subseteq T_{\mathcal{F}}$

Returns : $\mathbf{x}_{n+1} \in \mathcal{X}$, $T_{n+1} \subseteq T_{\mathcal{F}}$

- 1: Draw $\mathbf{x}^* \sim q_{\mathbf{x}_n}(\cdot)$
- 2: Solve $\tilde{\phi}_{T_n}(\mathbf{x}^*) : \mathbf{x}^* = \sum_{i=1}^n \alpha_i \mu(t_i)$
- 3: Draw a tree $t \in T_n$ ($\sim p_{\alpha}$), a context $c \in C_{\mathcal{F}}(t)$ ($\sim q_t$), and a subtree $\tau \in T_{\mathcal{F}} \sim c^{-1}\pi(\cdot)$
- 4: $t_{n+1} \leftarrow c[\tau]$, $T_{n+1} \leftarrow T_n \cup \{t_{n+1}\}$
- 5: Accept \mathbf{x}^* with probability

$$\alpha(\mathbf{x}_n, \mathbf{x}^*) = \min \left\{ 1, \frac{\hat{\phi}_{T_n}(\mathbf{x}^*) q_{\mathbf{x}^*}(\mathbf{x}_n)}{\hat{\phi}_{T_n}(\mathbf{x}_n) q_{\mathbf{x}_n}(\mathbf{x}^*)} \right\} = \min \left\{ 1, \frac{\tilde{\phi}_{T_n}(\mathbf{x}^*)}{\tilde{\phi}_{T_n}(\mathbf{x}_n)} \right\}$$

Theorem

The distribution of the \mathbf{x}_i generated by this algorithm converges to $\hat{\phi}_{T_{\mathcal{F}}}$.

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Setting

Let $\hat{t} \in T_{\mathcal{F}}$, $d(\cdot, \cdot)$ a distance between trees $\phi(\cdot) = \exp(-d(\hat{t}, \cdot))$.
Goal: Find t maximizing ϕ (i.e. find \hat{t}).

Weighted tree automaton $\mathcal{A} = \langle Q, \mathcal{F}, R \rangle$, where $Q = \{q_1, q_2\}$, $\mathcal{F} = \{f(\cdot, \cdot), a\}$, and R is the set of rules:

$$\begin{array}{ll} q_1 \xrightarrow{0.9} f(q_1, q_2) & q_2 \xrightarrow{0.4} f(q_2, q_2) \\ q_1 \xrightarrow{0.1} a & q_2 \xrightarrow{0.6} a \end{array}$$

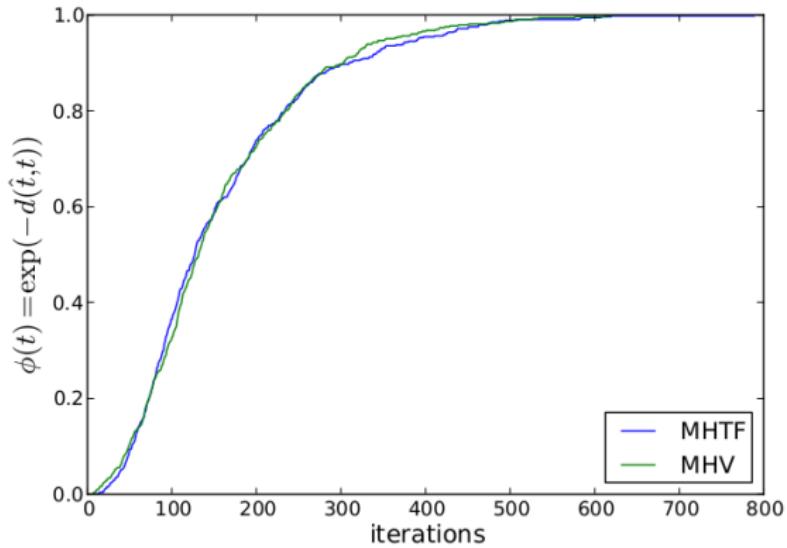
We run both algorithm

MHTF : MH in $T_{\mathcal{F}}$ using the stochastic series π induced by \mathcal{A} to generate the subtrees.

MHV : MH in the representation space induced by \mathcal{A} using π to generate the subtrees.

Results

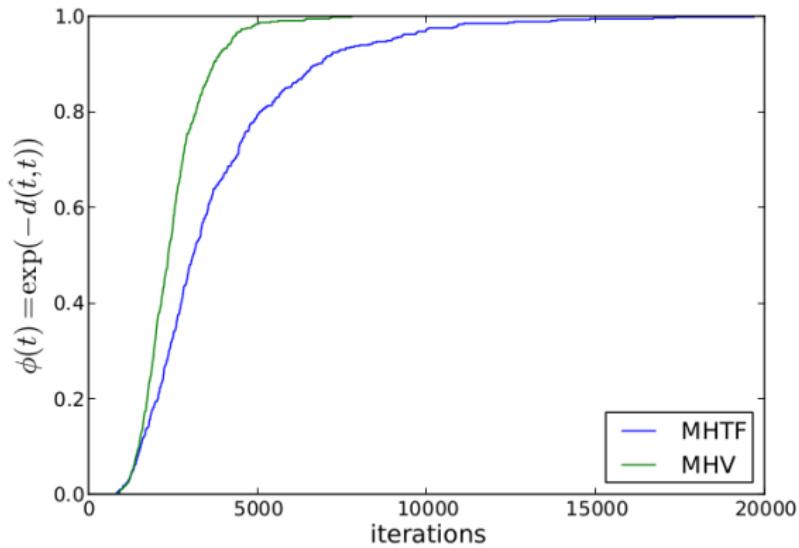
Size of target: 20 nodes



$$\hat{t} = f(f(f(f(f(f(f(a, a), a), a), a), a), a), f(a, a)), a)$$

Results

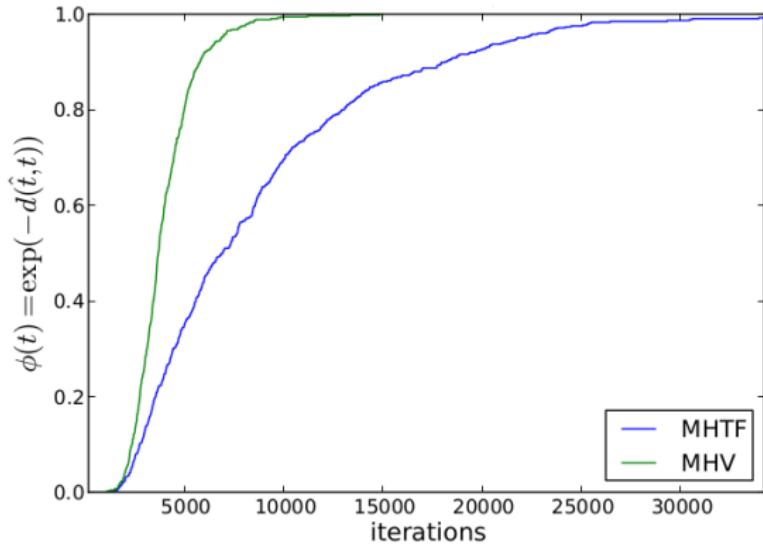
Size of target: 90 nodes



$$\begin{aligned}\hat{t} = & f(f(f(f(f(f(a, f(f(f(f(a, a), a), a), f(f(f(a, f(a, f(a, f(f(f(f(a, f(a, a)), a), a)))), a)))), a), a), a), a), \\& f(a, f(f(f(a, a), f(f(f(a, f(a, a)), a), a)))), a), a), a), f(f(f(a, a), f(f(a, a), f(a, \\& f(f(a, a), a)))), a), a), f(f(f(f(a, a), a), f(f(a, a), a))), a)\end{aligned}$$

Results

Size of target: 110 nodes



$$\begin{aligned}\hat{t} = & f(f(f(f(f(a, a), f(a, f(a, a))), f(f(a, a), f(a, a))), f(f(a, f(f(f(a, a), f(a, f(a, f(a, a))))), a)), f(f(f(f(a, f(f(a, a), f(a, f(f(f(a, a), a), f(a, a)))), f(f(f(a, a), a), f(a, a)), f(f(a, a), f(a, a))))))), a), f(f(a, a), \\& f(f(f(a, a), f(f(f(f(a, a), a), f(a, a)))), f(f(f(f(a, f(f(f(a, a), a), a)), f(f(a, a), a), f(a, a)), f(f(a, a), f(a, a))))))), a), f(f(a, a), \\& f(f(f(a, a), f(a, f(a, a)))), a))), f(a, f(a, f(a, a))), a)\end{aligned}$$

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Conclusion and Perspectives

- Contributions
 - ▶ Complexity study of problems involving a representation space
 - ▶ Algorithm to solve a tree problem in a representation space
 - ▶ Convergence guarantees
- Perspectives
 - ▶ Quantitative convergence bounds
 - ▶ How to choose the representation space?
 - ▶ Solve other tree problems in the representation space (e.g. classification)

-  Berstel, J. and Reutenauer, C. (1982).
Recognizable formal power series on trees.
Theoretical Computer Science, 18(2):115–148.
-  Blondel, V. and Canterini, V. (2003).
Undecidable problems for probabilistic automata of fixed dimension.
-  Fort, G., Moulines, E., and Priouret, P. (2012).
Convergence of adaptive and interacting Markov chain Monte Carlo algorithms.
-  Talton, J. O., Lou, Y., Lesser, S., Duke, J., Mech, R., and Koltun, V. (2011).
Metropolis procedural modeling.

Thank You.

Linear mapping μ

$$\begin{array}{ccc} \begin{array}{c} f \\ / \quad \backslash \\ a \quad a \end{array} & \Rightarrow & \begin{array}{c} \mu(f) \in \mathcal{L}(V \times V; V) \\ / \quad \backslash \\ V \ni \mu(a) \quad \mu(a) \in V \end{array} \\ \mathcal{F} = \{f(\cdot, \cdot), a\} & & \mu(t) = \mu(f)(\mu(a), \mu(a)) \end{array}$$

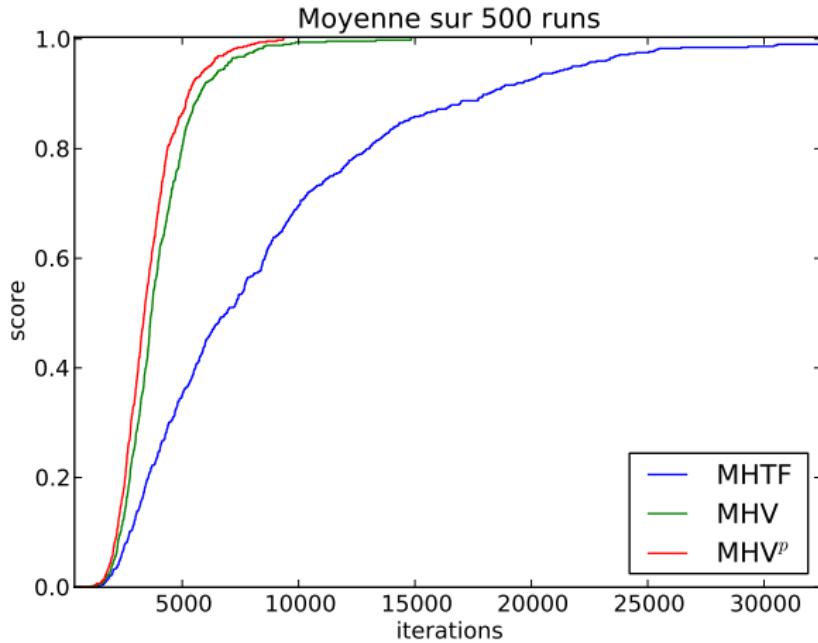
Vitesse d'exécution des algorithmes

- $\hat{t} = f(f(f(f(f(f(f(a, a), a), a), a), a), a), a)$

Algorithme	Nombre moyen d'itérations	Temps moyen d'exécution	Temps moyen par itération
MHTF	159	0.12 s	0.8 ms
MHV	165	0.17 s	1 ms

- $\hat{t} = f(f(f(f(f(f(f(a, f(f(f(f(f(a, f(f(f(f(f(a, f(f(f(f(f(a, f(a), a), a)$

Algorithme	Nombre moyen d'itérations	Temps moyen d'exécution	Temps moyen par itération
MHTF	3778	7.93 s	2.1 ms
MHV	2595	5.5 s	2.12 ms



$$\hat{t} = f(f(f(f(f(a, a), f(a, f(a, a))), f(f(a, a), f(a, a))), f(f(a, f(f(f(a, a), f(a, f(a, f(a, a))))), a)), f(f(f(f(a, f(f(a, a), f(a, f(f(f(a, a), a), f(a, a)))), f(f(f(a, a), a), f(f(a, a), f(f(a, a))))))), a), f(f(a, a), f(f(a, a), f(a, f(a, a)))), a))$$