Lecture 5: Exact inference

- Queries
- Inference in chains
- Variable elimination
- Without evidence
- With evidence
- Complexity of variable elimination

Queries

Bayesian networks can answer questions about the underlying

probability distribution:

- Likelihood: what is the probability of a given value assignment for a subset of variables Y?
- Conditional probability query: what is the probability of different variables Z? I.e. compute P(Y|Z = z)value assignments for query variables Y given evidence about
- Most probable evidence (MPE): given evidence Z = z, find an which has the highest probability: instantiation of all other variables in the network, W = X - Z,

$$MPE(W|Z = z) = \arg\max_{w} P(W = w|Z = z)$$

Queries (continued)

values to the variables in Y given that Z = z: given a subset of variables Y, find the most likely assignment of **Maximum a posteriori (MAP) query**: given evidence Z = z, and

$$MAP(Y|Z=z) = \arg\max_{y} P(Y=y|Z=z)$$

Examples of MAP queries:

In speech recognition, given a speech signal, one can attempt

generated the signal. to reconstruct the most likely sequence of words that could have

In classification, given the training data and a new example, we

want to determine the most probable class label of the new

example.

Complexity of inference

notes for details). whether P(X = x) > 0 is NP-hard (see Friedman and Koller's Given a Bayesian network and a random variable X, deciding

This implies that there is no general inference procedure that

will work efficiently for all network configurations

But for particular families of networks, inference can be done efficiently.

Likelihood inference in simple chains

Consider a simple chain of nodes:

$$A \to B \to C \to D$$

How do we compute P(B)?

$$P(B) = \sum_{a} P(A = a) P(B|A = a)$$

and B has m possible values, this requires O(km) operations: kAll the numbers required are in the CPTs. If A has k possible values multiplications and k-1 additions for each of the m values of B.

P(C)node C. Now how do we compute P(C)? We use P(B), which is already computed, and the local CPT of $= \sum P(B=b)(C|B=b)$ $\sum_{b} \left(\sum_{a} P(A = a) P(B = b | A = a) \right) P(C | B = b)$ Inference in simple chains (2) $A \to B \to C \to D$

Inference in simple chains (3)

$$X_1 \to X_2 \to \ldots \to X_n$$

How do we compute $P(X_n)$?

values of variable X), and the algorithm is linear in the number of We compute $P(X_2), \ldots P(X_n)$ iteratively. Each step only takes variables $O(|X_i| \cdot |X_{i+1}|)$ operations (where |X| is the number of possible

If we would have generated the whole joint distribution and summed

out, we would have needed $O((\max_i |X_i|)^n)$ operations!

Elimination of variables in chains

Suppose we want to compute P(D): Let us examine the chain example again: $A \to B \to C \to D$.

$$P(D) = \sum_{A,B,C} P(A,B,C,D)$$

$$= \sum_{A,B,C} P(A)P(B|A)P(C|B)P(D|C)$$

$$= \sum_{C} P(D|C) \sum_{B} P(C|B) \sum_{A} P(B|A)P(A)$$

can compute a factor $f_1(B)$, with one entry for each value of B. Then we can use this to compute a factor $f_2(C)$ etc. The innermost summation depends only on the value of B. So we

This is a form of dynamic programming

Pooling

Consider the case when a node has more than one parent, e.g.:



How do we compute P(C)?

$$P(C) = \sum_{A} P(C|A)P(A) = \sum_{A} P(C|A) \sum_{E,B} P(A|B, E)P(E)P(B)$$

A Bayes network is called a polytree if the underlying undirected

graph is a tree.

What if the network is not a polytree?



Suppose we want to compute P(W).

$$P(W) = \sum_{R,S,C} P(W,R,S,C) = \sum_{R,S,C} P(W|R,S)P(R|C)P(S|C)P(C)$$
$$= \sum_{R,S} P(W|R,S) \sum_{C} P(R|C)P(S|C)P(C)$$

Note that in this case we have a more complex factor, which

depends on two variables.

Variable elimination without evidence

Given: A Bayes network and a set of query variables $Y_1,\ldots Y_k$

1. Initialize the set of factors: $F = \{P(X_i | Parents(X_i))\}, \forall i$.

2. Let $\{Z_1 \dots Z_m\} = \{X_1, \dots, X_m\} - \{Y_1 \dots Y_k\}$

3. For i = 1 ... m do:

(a) Extract from F all factors $f_1, \ldots f_r$ mentioning Z_i

(b) Let $f' = \prod_{j=1}^r f_j$ (c) Let $f'' = \sum_{Z_i} f'$.

(d) Insert f'' in F

4. Return $\prod_{f \in F} f$

Steps (a) and (b) eliminate variable Z_i ; this is where the

computations actually take place



This example is taken from Koller and Friedman's notes:



Note that we can use any ordering of the variables during

elimination

Predictive inference with evidence in chains

Suppose we know that A = a. How do we compute P(C|A = a)?

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$$\begin{aligned} \langle C|A = a \rangle &= \frac{P(C, A = a)}{P(A = a)} = \frac{\sum_{B} P(C, B, A = a)}{P(A = a)} \\ &= \frac{\sum_{B} P(C|B) P(B|A = a) P(A = a)}{P(A = a)} \\ &= \sum_{D} P(C|B) P(B|A = a) \end{aligned}$$

Without knowing A, computing P(C, A) required another factor:

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$$P(C, A) = \sum_{B} P(C, B, A) = \sum_{B} P(C|B) \sum_{A} P(B|A)P(A)$$

$$P(C, A) = \sum_{B} P(C, B, A) = \sum_{B} P(C|B) \sum_{A} P(B|A)P(A)$$

instead of $\sum_{A} P(B|A)P(A)$. We eliminated the factor *inconsistent* Computing P(C, A = a) requires using P(B|A = a)P(A = a)

Causal inference with evidence in chains

that B = b. How do we compute P(A|B = b)? Again the chain example: $A \rightarrow B \rightarrow C \rightarrow D$. Suppose we know

We apply Bayes rule:

$$P(A|B = b) = \frac{P(A, B = b)}{P(B = b)}$$

We do not need to compute P(B = b), that comes out of summing

the numerators for all values of A.

P(A, B = b) can be computed using Bayes rule:

P(A, B = b) = P(B = b|A)P(A).

This can be viewed as a message passing from B to A.

Causal inference with evidence in chains (2)

$$A \to B \to C \to D$$

Suppose we know that C = c. How do we compute P(B|C = c)?

$$P(B|C = c) = \frac{P(B, C = c)}{P(C = c)} = \frac{P(C = c|B)P(B)}{P(C = c)}$$

computed using forward inference, just like before. P(C = c|B) is known from the CPT of node C. P(B) can be

from A (forward pass), and performs the computation. B receives some information from C (backward pass) and some

Inference with evidence in polytrees



How do we compute P(E|C = t)? We need P(E, C = t).

$$P(E, C = t) = \sum_{A,B} P(E, A, C = t, B)$$
$$= \sum_{A} P(C = t|A) \sum_{B} P(E)P(B)P(A|E, B)$$



Example: Asia network

Variable elimination with evidence

Given: A Bayes network, a set of query variables $Y_1, \ldots Y_k$, and evidence $u_1, \ldots u_l$

1. Initialize the set of factors: $F = \{P(X_i | Parents(X_i))\}, \forall i$.

2. For each factor, if it contains u_i , retain only the appropriate

portion (to be consistent with the evidence)

3. Let
$$\{Z_1 \dots Z_m\} = \{X_1, \dots, X_n\} - \{Y_1 \dots Y_k\} - \{U_1 \dots U_l\}$$

4. FUL $i = 1 \dots i l$ uu.

(a) Extract form F all factors $f_1, \ldots f_r$ mentioning Z_i

(b) Let
$$f' = \prod_{j=1}^r f_j$$

(c) Let $f'' = \sum_{Z_i} f'$.

(d) Insert
$$f''$$
 in F
5. Return $\prod f = f$

 $I I f \in F J$

Complexity of variable elimination

- factor fWe need at most O(n) multiplications to create one entry in a
- The size of a factor f containing m variables is k^m , where k is the maximum arity of a variable
- We need O(n) additions

So to be efficient, it is important to have small factors.

Induced graph

structure. We will try to understand the size of the factors in terms of the graph

appear in an intermediate factor f generated by variable elimination. $X_1 \dots X_n$ where X_i and X_j are connected by an edge if they both $Y_1, \ldots Y_k$, the **induced graph** H is an undirected graph over Given a Bayes net structure G and an elimination ordering

Example: Asia network

For our previous example, let us construct the induced graph:



graph.

vertices can be added)

each vertex is connected to every other vertex

- A clique is a maximal complete subgraph (one to which no
- - A complete subgraph of H is a subset of vertices such that



Cliques

Complexity of variable elimination

Theorem:

1. Every clique in the induced graph corresponds to an

intermediate factor in the computation

2. Every factor generated during variable elimination is a subset of

some maximal clique.

See Koller and Friedman notes for the proof details.

Therefore, complexity is exponential in the size of the largest clique

Consequence: Polytree inference

network (which includes all the CPTs). For the class of polytree networks, the problem of computing P(X, y) for any X can be solved in time linear in the size of the

clique is the largest family in the graph. The conclusion follows. induced graph is exactly the moral graph of the tree. So the largest **Proof:** We can order the nodes from the leaves inward. The

Heuristics for node ordering

- to <u>1</u>. Maximum cardinality: Number the nodes from 1 to n, always assigning the next number to the vertex having the largest set of previously numbered neighbors. Then eliminate nodes from $ar{n}$
- Minimum discrepancy: Always eliminate the node that causes the fewest edges to be added to the induced graph
- Minimum size/weight: Eliminate the node that causes the smallest clique to be created (either in terms of number of nodes, or in terms of number of entries).

Summary

- General exact inference in Bayesian networks is NP-hard
- Variable elimination is a general algorithm for exact inference
- By analyzing variable elimination we can see the "easy" cases for inference:
- when the net is a polytree
- when the maximum clique of the induced graph is small
- Heuristics for ordering work pretty well in practice